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**Airat Bikchentaev & Rinat Yakushev**

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# States on Symmetric Logics: Conditional Probability and Independence. II

Airat Bikchentaev · Rinat Yakushev

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**Abstract** We study the notions of conditional probabilities, independence and  $\varepsilon$ -independence for states on symmetric logics. We prove that a non-atomic state on the logic with the Lyapunov's property is determined by its specification of independent events. We present the examples of (1)  $\Delta$ -subadditive but is not subadditive and (2) two-valued non  $\Delta$ -subadditive states on symmetric logic. We investigate the independence relation transitivity for a  $\Delta$ -subadditive state.

We also study continuity properties of conditional probabilities and  $\varepsilon$ -independence relation with respect to natural pseudometric for  $\Delta$ -subadditive state. Finally, we pose two open problems.

**Keywords** Quantum logic · State · Conditional probability · Independence · Symmetric difference

## 1 Introduction and Preliminaries

Measure theory problems for quantum logics (particular, Boolean algebras and  $\sigma$ -algebras) of sets are an actual field of mathematical activity cf. [6, 7, 15, 16] and references therein. The notion of conditional probability is the principle instrument of the classical probability theory (cf., e.g., [5, Chap. V]).

This paper continues the first author's study begun in [2]; so we retain the notation and terminology used there. Our aim is to study the notions of conditional probabilities, independence and  $\varepsilon$ -independence for states on symmetric logics. We prove that a non-atomic state on the logic with the Lyapunov's property is determined by its specification of independent events.

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A. Bikchentaev (✉) · R. Yakushev  
N.I. Lobachevskii Institute of Mathematics and Mechanics, Kazan Federal University,  
Kremlevskaya 18, 420008 Kazan, Russian Federation  
e-mail: [Airat.Bikchentaev@kpfu.ru](mailto:Airat.Bikchentaev@kpfu.ru)

R. Yakushev  
e-mail: [Rinat.Yaqushev@kpfu.ru](mailto:Rinat.Yaqushev@kpfu.ru)

We present the examples of (1)  $\Delta$ -subadditive but not subadditive and (2) two-valued non  $\Delta$ -subadditive states on symmetric logics.

We investigate the independence relation transitivity for  $\Delta$ -subadditive states. We also study continuity properties of conditional probabilities and  $\varepsilon$ -independence relation with respect to natural pseudometric for  $\Delta$ -subadditive state. We prove that in this pseudometric space any “triangle” possesses a “perimeter” less than or equal to 2.

Finally, we pose two open problems.

Let us recall [6] that the set  $\mathcal{E}$  of subsets of  $\Omega$  is called a *concrete quantum logic* if the following conditions hold true:

- (1)  $\Omega \in \mathcal{E}$ ;
- (2)  $A \in \mathcal{E} \Rightarrow A^c = \Omega \setminus A \in \mathcal{E}$ ;
- (3)  $A, B \in \mathcal{E}, A \cap B = \emptyset \Rightarrow A \cup B \in \mathcal{E}$ .

Let  $\mathcal{S}(\Omega)$  be the set of all subsets of  $\Omega$ . Let us consider the following condition for  $\mathcal{E} \subset \mathcal{S}(\Omega)$ :

- (4)  $A, B \in \mathcal{E}, A \subset B \Rightarrow B \setminus A \in \mathcal{E}$ .

It seems clear that a family  $\mathcal{E} \subset \mathcal{S}(\Omega)$  is a concrete quantum logic if and only if conditions (1) and (4) hold.

We say that the mapping  $m : \mathcal{E} \rightarrow [0, 1]$  is a *state* (or a *probability measure*) on the concrete logic  $\mathcal{E}$ , if  $m(\Omega) = 1$  and  $m(A \cup B) = m(A) + m(B)$  for all  $A, B \in \mathcal{E}, A \cap B = \emptyset$ . Let us denote by  $P(\mathcal{E})$  the set of all states on logic  $\mathcal{E}$ . Recall that the state  $m \in P(\mathcal{E})$  is called *subadditive* ([15], p. 829) if for each  $A, B \in \mathcal{E}$  there exists a set  $C \in \mathcal{E}$  such that  $C \supset A \cup B$  and, moreover,  $m(C) \leq m(A) + m(B)$ .

In what follows the elements of  $\mathcal{E}$  will be called events. Every minimal element of  $\mathcal{E} \setminus \{\emptyset\}$  with respect to inclusion is called an *atom* in  $\mathcal{E}$ .

Let  $\nu : \mathcal{E} \rightarrow \mathbb{R}_+^n$  ( $n \geq 1$ ) be a (vector) measure ( $\nu(A \cup B) = \nu(A) + \nu(B)$  for all  $A, B \in \mathcal{E}, A \cap B = \emptyset$ ). An event  $A \in \mathcal{E}$  is  $\nu$ -*atom* if  $\nu(A) > 0$  and if for any event  $B \subset A$ , either  $\nu(B) = \nu(A)$  or  $\nu(B) = 0$ . A (vector) measure  $\nu$  is *nonatomic* if it has no  $\nu$ -atoms.

The set  $\mathcal{S}(\Omega)$  is a group with respect to the symmetric difference operation:  $A \Delta B = (A \setminus B) \cup (B \setminus A)$ . Since  $(X \Delta Y) \Delta Z = X \Delta (Y \Delta Z)$  for arbitrary  $X, Y, Z \in \mathcal{S}(\Omega)$  one can write  $X \Delta Y \Delta Z$  without brackets which we need to order operations. Thus

$$A^c \Delta B = (\Omega \Delta A) \Delta B = A \Delta (\Omega \Delta B) = \Omega \Delta (A \Delta B) = (A \Delta B)^c,$$

$$A^c \Delta B^c = (\Omega \Delta A) \Delta (\Omega \Delta B) = (\Omega \Delta A \Delta \Omega) \Delta B = ((\Omega \Delta \Omega) \Delta A) \Delta B = A \Delta B.$$

A concrete logic  $\mathcal{E}$  is said to be *symmetric* [12, Definition 3.2], if

- (5)  $A, B \in \mathcal{E} \Rightarrow A \Delta B \in \mathcal{E}$ .

These logics were investigated e.g. in [2, 4, 8, 9, 12, 13]. A family  $\mathcal{E} \subset \mathcal{S}(\Omega)$  is a symmetric logic if and only if conditions (1) and (5) hold [2, Proposition 1].

*Example 1.1* Let  $n \in \mathbb{N}$  and  $\Omega = \{1, 2, \dots, 2n\}$ . Then the family

$$\Omega_{\text{even}} = \{A \subset \Omega : \text{card}(A) = 2k, k = 0, 1, 2, \dots, n\}$$

is a symmetric logic on  $\Omega$ .

*Example 1.2* Let  $\mathcal{E} \subset \mathcal{S}(\Omega)$  be a concrete quantum logic and  $A \in \Omega, A \neq \emptyset$ . Then the family  $\mathcal{E}_A = \{B \in \mathcal{E} : B \subset A\}$  is a concrete quantum logic with the greatest element  $A$ . Moreover, if  $\mathcal{E}$  is a symmetric logic, then  $\mathcal{E}_A$  is also a symmetric logic.

## 2 Conditional Probability and Independence

In what follows  $\mathcal{E}$  will be a symmetric logic on  $\Omega$ . Let  $m \in P(\mathcal{E})$ . For  $A, B \in \mathcal{E}$  we define [2]

$$\tilde{m}(A, B) = \frac{m(A) + m(B) - m(A \Delta B)}{2}.$$

Thus  $\tilde{m}(A, A) = m(A)$ ,  $\tilde{m}(A, B) = \tilde{m}(B, A)$ ,  $\tilde{m}(A, B) + \tilde{m}(A^c, B) = m(B)$ , and  $\tilde{m}(A, B) = 0$  if  $A \cap B = \emptyset$ .

**Definition 2.1** [2] Let  $m \in P(\mathcal{E})$  and  $A, B \in \mathcal{E}$ . Let us say that the conditional probability of an event  $B$  under condition of another event  $A$  with  $m(A) > 0$  is the value

$$\frac{\tilde{m}(A, B)}{m(A)},$$

which will be denoted by  $m(B | A)$ .

Conditional probabilities may take negative values as well as positive ones. They are indefinite if  $m(A) = 0$ . If  $m$  is a probability on some Boolean algebra  $\mathcal{E}$ , then

$$\tilde{m}(A, B) = m(A \cap B), \quad m(B | A) = \frac{m(A \cap B)}{m(A)},$$

thus our definition coincides with the classical one.

**Definition 2.2** [2] Let  $m \in P(\mathcal{E})$  and  $A, B \in \mathcal{E}$ . Then two events  $A$  and  $B$  are independent if

$$\tilde{m}(A, B) = m(A)m(B).$$

It is proved in Theorem 1 of [2] that the following conditions are equivalent:

- (a) events  $A$  and  $B$  are independent;
- (b) events  $A$  and  $B^c$  are independent;
- (c) events  $A^c$  and  $B$  are independent;
- (d) events  $A^c$  and  $B^c$  are independent.

Since  $\tilde{m}(A, \Omega) = m(A) = m(A)m(\Omega)$ , events  $A$  and  $\Omega$  are independent. The events  $A$  and  $B$  are independent if and only if  $m(A | B) = m(A | B^c)$  [2, Theorem 2]. Two events  $A$  and  $B$  are independent if  $m(A | B) = m(A)$  [2, p. 103].

**Theorem 2.3** Let  $m \in P(\mathcal{E})$  and  $A, B \in \mathcal{E}$ ,  $0 < m(A) < 1$ . The following conditions are equivalent:

- (i)  $m(B | A) = m(B)$ ;
- (ii)  $m(B | A^c) = m(B)$ ;
- (iii)  $m(B^c | A) = m(B^c)$ ;
- (iv)  $m(B^c | A^c) = m(B^c)$ .

*Proof* All these conditions are equivalent to the independence of  $A$  (respectively  $A^c$ ) and  $B$  (respectively  $B^c$ ). □

**Proposition 2.4** Let  $A, B, C \in \mathcal{E}$  and  $C \subset B \subset A, m(B) > 0$ . Then

$$m(C | A) = m(C | B)m(B | A).$$

*Proof* As  $A, B, C$  are contained in a Boolean subalgebra of  $\mathcal{E}$ , the classical proof works.  $\square$

If  $m(A) > 0$  and  $m(B) > 0$ , then we obtain an analogue of *Bayes formula*:

$$m(A | B) = \frac{m(A)m(B | A)}{m(B)}.$$

If  $0 < m(A) < 1$ , then  $m(B) = m(B | A)m(A) + m(B | A^c)m(A^c)$ .

**Proposition 2.5** Let  $\mathcal{E}$  be a symmetric logic and  $m \in P(\mathcal{E})$ ,  $A, B \in \mathcal{E}$  and  $m(A), m(B) \in (0, 1)$ . The following conditions are equivalent:

- (i) events  $A, B$ , and  $A \Delta B$  are pairwise independent;
- (ii)  $m(A) = m(B) = m(A \Delta B) = 1/2$ .

*Proof* (i)  $\Rightarrow$  (ii). We have

$$m(A) + m(B) - m(A \Delta B) = 2m(A)m(B), \tag{1}$$

$$m(A) + m(A \Delta B) - m(B) = 2m(A)m(A \Delta B), \tag{2}$$

$$m(B) + m(A \Delta B) - m(A) = 2m(B)m(A \Delta B). \tag{3}$$

The sum of formulas (1) and (2) allows us to have the reduction  $m(B) + m(A \Delta B) = 1$ ; similarly combination of formulas (1) and (3) provides us with the relation  $m(A) + m(A \Delta B) = 1$ . Thus  $m(A) = m(B)$ . Again combination of formulas (2) and (3) gives us  $m(A) + m(B) = 1$ .

The implication (ii)  $\Rightarrow$  (i) can be verified by direct computation.  $\square$

**Proposition 2.6** If  $m$  is a nonatomic state on the symmetric logic  $\mathcal{E}$  and  $A \in \mathcal{E}, m(A) > 0$ , then there exist events  $B \subseteq A, C \subseteq A$  such that  $\tilde{m}(B, C) = m(B)m(C)$ .

*Proof* Coincides with the proof of [3, Lemma 1].  $\square$

**Proposition 2.7** Let  $\mathcal{E}$  be a symmetric logic and  $m \in P(\mathcal{E})$ . If  $A, B \in \mathcal{E}, A \cap B = \emptyset$  and  $m(A)m(B) > 0$ , then the events  $A$  and  $B$  are dependent.

*Proof* We have

$$0 = \frac{m(A) + m(B) - m(A \Delta B)}{2} \neq m(A)m(B). \tag{4} \quad \square$$

### 3 Independence and Determination of States

We say that  $m, \mu \in P(\mathcal{E})$  have identical independent events, if, for any pair of events  $A$  and  $B, \tilde{m}(A, B) = m(A)m(B)$  if and only if  $\tilde{\mu}(A, B) = \mu(A)\mu(B)$ .

**Definition 3.1** A symmetric logic  $\mathcal{E}$  has the Lyapunov's property (write  $\mathcal{E} \in (\text{LP})$ ) if for every pair  $m, \mu \in P(\mathcal{E})$  such that the vector measure  $\nu = (m, \mu)$  is nonatomic, the range of  $\nu$  is convex.

**Theorem 3.2** Let  $\mathcal{E} \in (\text{LP})$  and  $m, \mu \in P(\mathcal{E})$  be so that at least one of them is nonatomic. If they have identical independent event pairs, then they coincide.

**Corollary 3.3** Let  $\mathcal{E} \in (\text{LP})$  and  $m, \mu \in P(\mathcal{E})$  be so that at least one of them is nonatomic. If they have identical mutually favorable events in the following sense: for any pair of events  $A$  and  $B$ ,  $\tilde{m}(A, B) \geq m(A)m(B)$  if and only if  $\tilde{\mu}(A, B) \geq \mu(A)\mu(B)$ , then  $m$  and  $\mu$  coincide.

**Corollary 3.4** Let  $\mathcal{E} \in (\text{LP})$  and  $m, \mu \in P(\mathcal{E})$  be so that at least one of them is nonatomic. If

$$m(A) = 1/2 \iff \mu(A) = 1/2,$$

then  $m = \mu$ .

Proof of Theorem 3.2 coincides with the second proof of Theorem 1 from [3] (we use our Proposition 2.5). It suffices to consider 2 independent events. Proofs of Corollaries 3.3 and 3.4 coincide with the proofs of Corollaries 1 and 2 of [3], respectively. A wide class of quantum structures with the Lyapunov's property was considered in [1].

#### 4 $\Delta$ -Subadditive States on Symmetric Logics

Let us say that a state  $m \in P(\mathcal{E})$  is  $\Delta$ -subadditive [4] if

$$m(A\Delta B) \leq m(A) + m(B) \quad \text{for any pair } A, B \in \mathcal{E}.$$

In [2] this term was introduced without the prefix  $\Delta$ . But since the usual subadditivity differs from this notion we feel obliged to correct ourselves.

The set of all  $\Delta$ -subadditive states is convex. A state  $m \in P(\mathcal{E})$  is  $\Delta$ -subadditive if and only if conditional probabilities are non-negative on it.

Let  $\mathcal{E}$  be a Boolean algebra,  $m \in P(\mathcal{E})$  and  $A, B_1, B_2 \in \mathcal{E}$ . If  $B_1 \subset B_2$ , then

$$m(B_1 | A) \leq m(B_2 | A),$$

i.e. the conditional probability is monotonic. A state  $m$  on a symmetric logic  $\mathcal{E}$  is  $\Delta$ -subadditive if and only if the conditional probability is monotonic [2, Theorem 3].

It is proved in Lemma 1 of [2] that a state  $m \in P(\mathcal{E})$  is  $\Delta$ -subadditive if and only if

$$m(A\Delta B) \leq m(A\Delta C) + m(C\Delta B) \quad \text{for all } A, B, C \in \mathcal{E}. \tag{4}$$

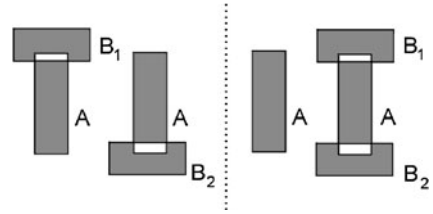
For a  $\Delta$ -subadditive state  $m \in P(\mathcal{E})$  via (4) we get

$$\tilde{m}(A, B) + \tilde{m}(A, C) - \tilde{m}(B, C) \leq m(A) \quad \text{for all } A, B, C \in \mathcal{E}.$$

*Example 4.1* Let  $\Omega = \{1, 2, 3, 4\}$ . Let us define the two-valued state  $m$  on symmetric logic  $\Omega_{\text{even}}$  by its values on atoms as follows:

$$m(\{1, 2\}) = m(\{1, 3\}) = m(\{1, 4\}) = 0.$$

**Fig. 1** Scheme of the sets  $A \Delta B_1$ ,  $A \Delta B_2$  and  $A$ ,  $A \Delta (B_1 \cup B_2)$



Then  $m$  is not  $\Delta$ -subadditive:

$$1 = m(\{3, 4\}) = m(\{1, 3\} \Delta \{1, 4\}) > m(\{1, 3\}) + m(\{1, 4\}) = 0.$$

The space of  $\Delta$ -subadditive states is a tetrahedron (a convex combination of 4 extreme states), the conditional probabilities achieve values between 0 and 1 exactly on this set.

Let  $\mathcal{E}$  be a Boolean algebra,  $m \in P(\mathcal{E})$  and  $A, B_1, B_2 \in \mathcal{E}$ . If the event  $A$  does not depend on the events  $B_1, B_2$  and  $B_1 \cap B_2 = \emptyset$ , then the events  $A$  and  $B_1 \cup B_2$  are independent.

**Theorem 4.2** *Let  $\mathcal{E}$  be a symmetric logic of subsets of a set  $\Omega$ ,  $\mathcal{A}$  be a Boolean algebra of subsets of  $\Omega$ ,  $\mathcal{E} \subset \mathcal{A}$  and  $A, B_1, B_2 \in \mathcal{E}$ . Let a state  $m$  on  $\mathcal{E}$  allow for an extension,  $\mathbf{m}$ , over  $\mathcal{A}$  as a signed measure. If the event  $A$  does not depend on the events  $B_1, B_2$  and  $B_1 \cap B_2 = \emptyset$ , then the events  $A$  and  $B_1 \cup B_2$  are independent.*

*Proof* Let  $B = B_1 \cup B_2$ . Since

$$m(A) + m(B_1) - m(A \Delta B_1) = 2m(A)m(B_1),$$

$$m(A) + m(B_2) - m(A \Delta B_2) = 2m(A)m(B_2),$$

we have

$$2m(A) + m(B) - m(A \Delta B_1) - m(A \Delta B_2) = 2m(A)m(B). \tag{5}$$

Since (see Fig. 1)

$$\begin{aligned} & m(A \Delta B_1) + m(A \Delta B_2) \\ &= m(A) + m(B_1) - 2\mathbf{m}(A \cap B_1) + m(A) + m(B_2) - 2\mathbf{m}(A \cap B_2) \\ &= m(A) + m(A \Delta B), \end{aligned}$$

we have via (5) the relation  $m(A) + m(B) - m(A \Delta B) = 2m(A)m(B)$ . This completes the proof.  $\square$

Let  $n \in \mathbb{N}$  and  $\Omega = \{1, 2, \dots, 2n\}$ . By Theorem 2.1 [4] every  $m \in P(\Omega_{\text{even}})$  can be extended to a signed measure  $\mathbf{m}$  over  $\mathcal{S}(\Omega)$ , the (Boolean) power algebra of  $\Omega$ . Thus Theorem 4.2 holds for any state  $m$  on  $\Omega_{\text{even}}$ .

**Example 4.3** Let  $\Omega = \{0, 1, 2, 3, 4, 5\}$  and the symmetric logic  $\mathcal{E}$  contain the sets

$$\begin{aligned} A &= \{1, 2, 3\}, & B_1 &= \{0, 1\}, & B_2 &= \{3, 4\}, & C &= \{2, 5\}, \\ D &= \{0, 2, 3\}, & E &= \{1, 2, 4\}, & F &= \{0, 2, 4\}, & \Omega & \end{aligned}$$



and their complements. The logic  $\mathcal{E}$  has 16 elements. Let  $m$  be a  $\Delta$ -subadditive state on  $\mathcal{E}$  such that

$$m(A) = m(D) = m(E) = 1/2, \quad m(B_1) = m(B_2) = 1/3, \quad m(F) = 3/8.$$

Then the event  $A$  does not depend on the events  $B_1, B_2$ , but the events  $A$  and  $B_1 \cup B_2$  are not independent. Since  $\{0\} = E^c \cap F = A^c \cap D$ ,  $\{1\} = A \cap B_1$ ,  $\{2\} = A \cap C$ ,  $\{3\} = B_2 \cap F^c$ ,  $\{4\} = A^c \cap B_2$ , and  $\{5\} = C \cap D^c$ , the Boolean algebra generated by  $\mathcal{E}$  is  $\mathcal{S}(\Omega)$ . Note that  $m$  cannot be extended as a signed measure on  $\mathcal{S}(\Omega)$ .

*Example 4.4* Let  $\Omega = \{1, 2, 3, 4, 5, 6\}$ . Again we define the state  $m$  on the symmetric logic  $\Omega_{\text{even}}$  by its values on atoms as follows:

$$\begin{aligned} m(\{1, 2\}) &= \frac{2}{3}, & m(\{1, 3\}) &= m(\{1, 4\}) = \frac{2}{9}, & m(\{1, 5\}) &= \frac{1}{3}, \\ m(\{1, 6\}) &= m(\{2, 3\}) = m(\{2, 4\}) = \frac{4}{9}, & m(\{2, 5\}) &= \frac{5}{9}, \\ m(\{2, 6\}) &= \frac{2}{3}, & m(\{3, 4\}) &= 0, & m(\{3, 5\}) &= \frac{1}{9}, & m(\{3, 6\}) &= \frac{2}{9}, \\ m(\{4, 5\}) &= \frac{1}{9}, & m(\{4, 6\}) &= \frac{2}{9}, & m(\{5, 6\}) &= \frac{1}{3}. \end{aligned}$$

Then  $m$  is  $\Delta$ -subadditive.

If  $\mathcal{E}$  is a symmetric logic, then every subadditive state  $m \in P(\mathcal{E})$  is  $\Delta$ -subadditive (hint:  $C \supset A \cup B \supset A \Delta B$ ), but the reverse implication is not true in general. The state  $m$  from Example 4.4 is  $\Delta$ -subadditive and is not subadditive: for  $A = \{3, 4\}$ ,  $B = \{1, 2, 4, 5\}$  we have the unique set  $C = \{1, 2, 3, 4, 5, 6\}$  such that  $C \supset A \cup B$  and  $1 = m(C) > m(A) + m(B) = 0 + 7/9 = 7/9$ .

**Theorem 4.5** Let  $m \in P(\mathcal{E})$ . The following conditions are equivalent:

- (i)  $m$  is  $\Delta$ -subadditive;
- (ii)  $\tilde{m}(A, B) \leq \min\{m(A), m(B)\}$  for all  $A, B \in \mathcal{E}$ .

*Proof* (i)  $\Rightarrow$  (ii). We have  $B = (B \Delta A) \Delta A$ , therefore  $m(B) \leq m(A \Delta B) + m(A)$ , i.e.  $m(B) - m(A \Delta B) \leq m(A)$ . The latter inequality is equivalent to the inequality  $\tilde{m}(A, B) \leq m(A)$ .

(ii)  $\Rightarrow$  (i). We have  $m(B) - m(A) - m(A \Delta B) \leq 0$  for all  $A, B \in \mathcal{E}$ . We replace  $B$  by  $B^c$ ; then  $m(B^c) - m(A) - m(A \Delta B^c) \leq 0$ , i.e.  $1 - m(B) - m(A) - 1 + m(A \Delta B) \leq 0$ . This completes the proof.  $\square$

**Theorem 4.6** Let  $m \in P(\mathcal{E})$  be  $\Delta$ -subadditive and  $A, B, C \in \mathcal{E}$ . Then  $|\tilde{m}(A, B) - \tilde{m}(A, C)| \leq m(B \Delta C)$ .

*Proof* We have  $B = (B \Delta C) \Delta C$  and  $C = (B \Delta C) \Delta B$ , therefore  $m(B) \leq m(B \Delta C) + m(C)$  and  $m(C) \leq m(B \Delta C) + m(B)$ , i.e.

$$|m(B) - m(C)| \leq m(B \Delta C). \tag{6}$$

By (4) we get  $m(A \Delta B) \leq m(A \Delta C) + m(C \Delta B)$  and  $m(A \Delta C) \leq m(A \Delta B) + m(B \Delta C)$ , therefore

$$|m(A \Delta C) - m(A \Delta B)| \leq m(B \Delta C). \tag{7}$$

Now we have via triangle inequality and formulas (6), (7) the relations

$$\begin{aligned} |\tilde{m}(A, B) - \tilde{m}(A, C)| &= \frac{|m(B) - m(C) - m(A \Delta B) + m(A \Delta C)|}{2} \\ &\leq \frac{|m(B) - m(C)|}{2} + \frac{|m(A \Delta C) - m(A \Delta B)|}{2} \\ &\leq m(B \Delta C). \end{aligned} \quad \square$$

**Theorem 4.7** Let  $m \in P(\mathcal{E})$  be  $\Delta$ -subadditive and  $A_1, A_2, \dots$  and  $B_1, B_2, \dots$  be two sequences of events, where  $m(B_n) \rightarrow 1$  as  $n \rightarrow \infty$ . Then

- (I)  $\lim_{n \rightarrow \infty} m(A_n) = \lim_{n \rightarrow \infty} \tilde{m}(A_n, B_n)$  under the following condition: there exists at least one of indicated limits;
- (II) if

$$\liminf_{n \rightarrow \infty} m(A_n) \geq a > 0, \tag{8}$$

then

$$\lim_{n \rightarrow \infty} \frac{m(A_n)}{\tilde{m}(A_n, B_n)} = 1.$$

*Proof* (I) We have

$$m(A_n) = \tilde{m}(A_n, B_n) + \tilde{m}(A_n, B_n^c) \quad \text{for all } n \in \mathbb{N}. \tag{9}$$

By Theorem 4.5,  $0 \leq \tilde{m}(A_n, B_n^c) \leq m(B_n^c) \rightarrow 0$  as  $n \rightarrow \infty$ . Hence  $\tilde{m}(A_n, B_n^c) \rightarrow 0$  as  $n \rightarrow \infty$  and  $\lim_{n \rightarrow \infty} m(A_n) = \lim_{n \rightarrow \infty} \tilde{m}(A_n, B_n)$ .

(II) Let  $i \in \mathbb{N}$  be so that  $m(A_n) \geq a/2$  for all  $n \geq i$ , and let  $j \in \mathbb{N}$  be so that  $m(B_n) \geq 1 - a/4$  for all  $n \geq j$ . We have for  $n \geq \max\{i, j\}$

$$\tilde{m}(A_n, B_n) = \frac{m(A_n) + m(B_n) - m(A_n \Delta B_n)}{2} \geq \frac{a/2 + 1 - a/4 - 1}{2} = \frac{a}{8}.$$

Then for  $n \geq \max\{i, j\}$  via (9) and Theorem 4.5 the relations

$$1 = \frac{m(A_n)}{m(A_n)} \leq \frac{m(A_n)}{\tilde{m}(A_n, B_n)} = \frac{\tilde{m}(A_n, B_n) + \tilde{m}(A_n, B_n^c)}{\tilde{m}(A_n, B_n)} = 1 + \frac{\tilde{m}(A_n, B_n^c)}{\tilde{m}(A_n, B_n)} \leq 1 + \frac{8}{a} m(B_n^c)$$

hold. Since  $m(B_n^c) \rightarrow 0$  as  $n \rightarrow \infty$ , this completes the proof. □

*Example 4.8* Condition (8), in general, cannot be omitted. Consider  $\Omega = [0, 1]$ , Borel  $\sigma$ -algebra  $\mathcal{E} \subset \mathcal{S}(\Omega)$  with Lebesgue measure  $m$  and the events  $A_n = [0, 1/n]$ ,  $B_n = [1/(2n), 1]$ .

### 5 On Independent Events

The notion of independence is fundamental in probability theory.

**Theorem 5.1** *Let  $m \in P(\mathcal{E})$  be  $\Delta$ -subadditive. The following conditions are equivalent:*

- (i) *all the events from  $\mathcal{E}$  are mutually independent;*
- (ii)  *$m(A) \in \{0, 1\}$  for all  $A \in \mathcal{E}$ ;*
- (iii) *the independence relation on  $\mathcal{E}$  is transitive.*

*Proof* (i)  $\Rightarrow$  (ii). Assume that there exists  $A \in \mathcal{E}$  with  $0 < m(A) < 1$ . Then we have  $0 < m(A^c) < 1$  and

$$\tilde{m}(A, A^c) = \frac{m(A) + m(A^c) - m(A\Delta A^c)}{2} = 0 \neq m(A)m(A^c). \tag{10}$$

(ii)  $\Rightarrow$  (i). Step 1. Let  $m(A) = m(B) = 0$ . Then

$$0 \leq m(A\Delta B) \leq m(A) + m(B) = 0 \tag{11}$$

and  $\tilde{m}(A, B) = 0 = m(A)m(B)$ .

Step 2. Let  $m(A) = 0, m(B) = 1$ . Then  $m(B^c) = 0$  and by Step 1 the events  $A$  and  $B^c$  are independent. Therefore, the events  $A$  and  $B$  are independent via Theorem 1 of [2].

Step 3. Let  $m(A) = m(B) = 1$ . Then  $m(A^c) = m(B^c) = 0$  and by Step 1 the events  $A^c$  and  $B^c$  are independent. Therefore, the events  $A$  and  $B$  are independent via Theorem 1 of [2].

The implication (i)  $\Rightarrow$  (iii) is obvious.

(iii)  $\Rightarrow$  (ii). Assume that the independence relation on  $\mathcal{E}$  is transitive and there exists  $A \in \mathcal{E}$  with  $0 < m(A) < 1$ . Then (a)  $A$  is independent of  $\emptyset$ ; (b)  $\emptyset$  is independent from  $A^c$ , but  $A$  dependent from  $A^c$  by (10). We obtain a contradiction. This completes the proof.  $\square$

**Proposition 5.2** *Let  $\mathcal{E}$  be a symmetric logic and  $m \in P(\mathcal{E})$  be  $\Delta$ -subadditive. Then the family  $\mathcal{E}_m = \{A \in \mathcal{E} : m(A) \in \{0, 1\}\}$  also form a symmetric logic.*

*Proof* Since  $\Omega \in \mathcal{E}_m$ , it suffices to show that  $A\Delta B \in \mathcal{E}_m$  for all  $A, B \in \mathcal{E}_m$ .

If  $m(A) = m(B) = 0$ , then  $m(A\Delta B) = 0$  by (11) and  $A\Delta B \in \mathcal{E}_m$ .

If  $m(A) = m(B) = 1$ , then  $m(A^c) = m(B^c) = 0$  and  $m(A^c\Delta B^c) = 0$  by (10). Since  $A^c\Delta B^c = A\Delta B$ , we have  $A\Delta B \in \mathcal{E}_m$ .

If  $m(A) = 1$  and  $m(B) = 0$ , then  $m(A^c) = 0$  and  $m(A\Delta B) = 1 - m(A^c\Delta B) = 1$ .  $\square$

From Theorem 5.1 and Proposition 5.2 we have

**Corollary 5.3** *Let  $\mathcal{E}$  be a symmetric logic and  $m \in P(\mathcal{E})$  be  $\Delta$ -subadditive. Then all events from  $\mathcal{E}_m$  are mutually independent and the independence relation on  $\mathcal{E}_m$  is transitive.*

### 6 On $\varepsilon$ -Independent Events

**Definition 6.1** Let  $m \in P(\mathcal{E})$  and  $A, B \in \mathcal{E}$ . For  $\varepsilon > 0$  two events  $A$  and  $B$  are  $\varepsilon$ -independent, if

$$|\tilde{m}(A, B) - m(A)m(B)| \leq \varepsilon.$$

**Proposition 6.2** *If an event  $A$  is  $\varepsilon$ -independent with itself for  $0 < \varepsilon \leq 1/4$ , then either  $m(A) \leq 2\varepsilon$  or  $m(A) \geq 1 - 2\varepsilon$ .*

*Proof* We have  $\tilde{m}(A, A) = m(A)$  and

$$|\tilde{m}(A, A) - m(A)m(A)| = |m(A) - m(A)m(A)| = |m(A)(1 - m(A))| \leq \varepsilon.$$

By solving this quadratic inequality with respect to  $m(A)$ , we get

$$m(A) \leq \frac{1 - \sqrt{1 - 4\varepsilon}}{2} \quad \text{or} \quad m(A) \geq \frac{1 + \sqrt{1 - 4\varepsilon}}{2}.$$

Finally, we apply the following inequalities for  $0 < \varepsilon \leq 1/4$ :

$$\frac{1 - \sqrt{1 - 4\varepsilon}}{2} \leq 2\varepsilon, \quad \frac{1 + \sqrt{1 - 4\varepsilon}}{2} \geq 1 - 2\varepsilon. \quad \square$$

**Theorem 6.3** *Let  $m \in P(\mathcal{E})$  and  $A, B \in \mathcal{E}$ . The following conditions are equivalent:*

- (i) *events  $A$  and  $B$  are  $\varepsilon$ -independent;*
- (ii) *events  $A$  and  $B^c$  are  $\varepsilon$ -independent;*
- (iii) *events  $A^c$  and  $B$  are  $\varepsilon$ -independent;*
- (iv) *events  $A^c$  and  $B^c$  are  $\varepsilon$ -independent.*

*Proof* Since  $X^{cc} = X$  ( $X \in \mathcal{S}(\Omega)$ ), it suffices to show that (i)  $\Rightarrow$  (ii). The following relations prove it:

$$\begin{aligned} & |\tilde{m}(A^c, B) - m(A^c)m(B)| \\ &= \left| \frac{1 - m(A) + m(B) - 1 + m(A\Delta B)}{2} - (1 - m(A))m(B) \right| \\ &= |-\tilde{m}(A, B) + m(A)m(B)| \leq \varepsilon. \quad \square \end{aligned}$$

It seems useful to compare Theorems 2.3 and 6.3 with Theorem 1 of [2]. Also the following assertion can be compared with Proposition 2.7.

**Proposition 6.4** *Let  $m \in P(\mathcal{E})$  be  $\Delta$ -subadditive and  $A \in \mathcal{E}$  be so that  $m(A) \leq \varepsilon$  or  $m(A) \geq 1 - \varepsilon$ . Then  $A$  and an arbitrary event  $B$  are  $\varepsilon$ -independent.*

*Proof* If  $m(A) \leq \varepsilon$ , then  $m(A)m(B) \leq \varepsilon$  and by Theorem 4.5 we have also  $\tilde{m}(A, B) \leq m(A) \leq \varepsilon$ . Therefore,

$$|\tilde{m}(A, B) - m(A)m(B)| \leq \max\{\tilde{m}(A, B), m(A)m(B)\} \leq \varepsilon.$$

If  $m(A) \geq 1 - \varepsilon$ , then  $m(A^c) \leq \varepsilon$  and, consequently,  $A^c$  and an arbitrary event  $B$  are  $\varepsilon$ -independent. By Theorem 6.3,  $A$  and an arbitrary event  $B$  are also  $\varepsilon$ -independent. This completes the proof.  $\square$

If  $m \in P(\mathcal{E})$  is  $\Delta$ -subadditive, then the mapping  $d_m : \mathcal{E} \times \mathcal{E} \rightarrow [0, \infty)$ , defined by the formula  $d_m(A, B) = m(A\Delta B)$  ( $A, B \in \mathcal{E}$ ) is a pseudometric on  $\mathcal{E}$  [2, Theorem 4]. Via (2) a state  $m$  is uniform continuous on  $\langle \mathcal{E}, d_m \rangle$ . Theorem 4.6 shows us that for every fixed  $A \in \mathcal{E}$  the mapping  $X \mapsto \tilde{m}(A, X)$  ( $X \in \mathcal{E}$ ) is uniform continuous on  $\langle \mathcal{E}, d_m \rangle$ .

**Theorem 6.5** Let  $m \in P(\mathcal{E})$  be  $\Delta$ -subadditive,  $A, A_n, B, B_n \in \mathcal{E}$  ( $n \in \mathbb{N}$ ). If  $A_n \rightarrow A, B_n \rightarrow B$  ( $n \rightarrow \infty$ ) on  $\langle \mathcal{E}, d_m \rangle$  and the events  $A_n$  and  $B_n$  are  $\varepsilon$ -independent for any  $n \in \mathbb{N}$ , then the events  $A$  and  $B$  are  $\varepsilon$ -independent.

*Proof* It follows from Theorem 4 of [2] that  $m(A_n) \rightarrow m(A)$  and  $m(B_n) \rightarrow m(B)$  as  $n \rightarrow \infty$ . Via Theorem 5 of [2] we have  $m(A_n \Delta B_n) \rightarrow m(A \Delta B)$  as  $n \rightarrow \infty$ . Therefore  $\tilde{m}(A_n, B_n) \rightarrow \tilde{m}(A, B)$  as  $n \rightarrow \infty$  and the inequality

$$|\tilde{m}(A, B) - m(A)m(B)| \leq \varepsilon$$

holds. This completes the proof. □

**Proposition 6.6** Let  $m \in P(\mathcal{E})$  be  $\Delta$ -subadditive. Then

$$d_m(A, B) + d_m(B, C) + d_m(C, A) \leq 2 \quad \text{for all } A, B, C \in \mathcal{E}. \tag{12}$$

*Proof* We have

$$m(A \Delta B) = 1 - m(A \Delta B^c),$$

$$m(B \Delta C) = 1 - m(B^c \Delta C),$$

$$m(C \Delta A) = 1 - m(A^c \Delta C).$$

Thus  $d_m(A, B) + d_m(B, C) + d_m(C, A) = 3 - m(A \Delta B^c) - m(B^c \Delta C) - m(A^c \Delta C)$ . Since

$$1 = m(A^c \Delta A) \leq m(A^c \Delta C) + m(C \Delta A) \leq m(A^c \Delta C) + m(C \Delta B^c) + m(B^c \Delta A),$$

we have (12). □

**Corollary 6.7** Let  $\mathcal{E}$  be a symmetric logic and  $m \in P(\mathcal{E})$  be  $\Delta$ -subadditive. Then

$$\tilde{m}(A, B) + \tilde{m}(B, C) + \tilde{m}(C, A) \geq m(A) + m(B) + m(C) - 1 \quad \text{for all } A, B, C \in \mathcal{E}.$$

### 7 Open Problems

If  $\mathcal{E}$  is a Boolean algebra then any state  $m \in P(\mathcal{E})$  is subadditive. There exists a concrete quantum logic which is not a Boolean algebra and all of its states are subadditive. The result was established in [14] with substantial help from the techniques developed in [10] and [11] (see also [15], p. 831).

**Problem 7.1** Let  $\mathcal{E}$  be a symmetric logic such that any state  $m \in P(\mathcal{E})$  is  $\Delta$ -subadditive. Is it true that  $\mathcal{E}$  is a Boolean algebra?

If  $\mathcal{E}$  is a  $\sigma$ -algebra then the pseudometric space  $\langle \mathcal{E}, d_m \rangle$  is complete for any state  $m \in P(\mathcal{E})$ .

**Problem 7.2** Let  $\mathcal{E}$  be both a symmetric logic and a  $\sigma$ -class,  $m \in P(\mathcal{E})$  be  $\Delta$ -subadditive. Is it true that the pseudometric space  $\langle \mathcal{E}, d_m \rangle$  is complete?

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