

# Generalization of the Newton Method for One Class of Nonconvex Mathematical Programming Problems

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This paper continues the study commenced in [1] and [2], where we consider iterative procedures, finding stationary points of smooth functions on a class of nonconvex sets. We generalize the Newton method applied for the solution of convex programming problems for the case, when constraints are represented as a set-theoretical difference of a convex set and a union of several convex sets. We formulate and prove a proposition on the convergence of the algorithm.

## 1. PROBLEM DEFINITION AND THE ALGORITHM

Consider the following problem: find a point which satisfies the necessary condition for a local minimum of a function  $\varphi(x)$  on a set  $X$  in the  $n$ -dimensional Euclidean space  $E^n$ , where  $X$  is a set-theoretical difference of certain sets  $F$  and  $\bigcup_{i=1}^l \text{int } G_i$ ; here  $F$  and  $G_i$ ,  $i = \overline{1, l}$ , are convex and closed, the sets of inner points of  $X$  and  $G_i$ ,  $i = \overline{1, l}$ , are nonempty;  $\varphi(x)$  belongs to the class  $C^2(Y)$  and is strongly convex on a certain convex set  $Y$  which includes  $X$ . Let each set  $G_i$ ,  $i = \overline{1, l}$ , at any its boundary point  $x$  have a unique supporting hyperplane, whose normal is assumed to be outer. This means that for all  $y \in G_i$  the unit normal vector  $n^i(x)$  satisfies the condition  $\langle n^i(x), y - x \rangle \leq 0$ . We also assume that with each  $i$  the unit normal vector  $n^i(x)$  is a continuous vector function at the boundary  $\partial G_i$  of the set  $G_i$ . The latter means that for any point  $x_* \in \partial G_i$  and an arbitrary sequence  $\{x_k\}$  which belongs to  $\partial G_i$  and converges to  $x_*$ , with each  $\varepsilon > 0$  one can find a number  $k(\varepsilon) \in N$  such that all  $k \geq k(\varepsilon)$  meet the inequality  $\|n^i(x_k) - n^i(x_*)\| < \varepsilon$ .

Below we use the following denotations:  $s^i(x)$  is a projection of a point  $x$  onto the set  $G_i$ ,  $n^i(x)$  is the unit normal vector of the hyperplane supporting to  $G_i$  at the point  $s^i(x)$ ,  $\Gamma^i(x) = \{e \in E^n : \langle n^i(x), e - s^i(x) \rangle \geq 0\}$ ,  $P(x) = F \cap \Gamma^1(x) \cap \Gamma^2(x) \cap \dots \cap \Gamma^l(x)$ . Projections  $s^i(x)$  are defined uniquely, because  $G_i$ ,  $i = \overline{1, l}$ , are convex sets of the Euclidean space  $E^n$ . Since each  $G_i$ ,  $i = \overline{1, l}$ , at any its boundary point  $x$  has only one supporting hyperplane, vectors  $n^i(x)$  and, therefore, half-spaces  $\Gamma^i(x)$ ,  $i = 1, 2, \dots, l$ , are uniquely defined for any  $x \in X$ . If with certain  $i$  points  $x$  and  $s^i(x)$  do not coincide, then vectors  $n^i(x)$  and  $x - s^i(x)$  have the same direction. Consequently,  $\langle n^i(x), x - s^i(x) \rangle > 0$ , i. e.,  $x \in \text{int } \Gamma^i(x)$ . But if  $x$  and  $s^i(x)$  coincide, then  $x$  belongs to the boundary of  $\Gamma^i(x)$ . Since  $x \in X \subset F$  and  $x \in \Gamma^i(x)$  with each  $i = 1, 2, \dots, l$ , we conclude that always  $x \in P(x)$ .

We propose to solve the stated problem by the following algorithm, constructing successive approximations.

**Step 0.** Put  $k = 0$ .

**Step 1.** Let  $x_k \in X$  be the  $k$ th approximation.

**Step 2.** Define points  $s^i(x_k)$ ,  $i = \overline{1, l}$ .

**Step 3.** Construct half-spaces  $\Gamma^i(x_k)$ ,  $i = \overline{1, l}$ .

**Step 4.** Construct the set  $P(x_k)$ .