

## T-Irreducible Extension of Polygonal Digraphs

A. V. Gavrikov<sup>1\*</sup>

<sup>1</sup>Saratov State University,  
ul. Astrahanskaya 83, Saratov, 410012 Russia

Received July 10, 2014

**Abstract**—Directed graphs are mathematical models of discrete systems. T-irreducible extensions are widely used in cryptography and diagnosis of discrete systems. A polygonal origraph is a digraph obtained from a circuit by some orientation of its edges. We propose an algorithm to construct a T-irreducible extension of a polygonal graph.

**DOI:** 10.3103/S1066369X16020031

**Keywords:** *polygonal graph, fault-tolerance of discrete systems, T-irreducible extension.*

### INTRODUCTION

We say that a *directed graph (digraph)* is a pair  $\vec{G} = (V, \alpha)$  consisting of nonempty digraph vertex set  $V$  and a digraph edge defining relation  $\alpha$  on the set  $V$ . A digraph edge is *incident* to the vertex  $v$  if  $v$  is either the starting or ending point of this edge. An *inclusion* of digraph  $\vec{G} = (V, \alpha)$  into the digraph  $\vec{H} = (W, \beta)$  is a bijective map  $\phi : V \rightarrow W$  such that  $((u, v) \in \alpha \rightarrow (\phi(u), \phi(v)) \in \beta) \quad \forall u, v \in V$ . In this case we say that the digraph  $\vec{G}$  can be included into the digraph  $\vec{H}$ . A *part* of digraph  $\vec{G} = (V, \alpha)$  is a digraph  $\vec{H} = (W, \beta)$  such that  $W \subseteq V$  and  $\beta \subseteq (W \times W) \cap \alpha$ . A part of digraph  $\vec{G} = (V, \alpha)$  is a *subgraph* of digraph  $\vec{H} = (W, \beta)$  if  $\beta = (W \times W) \cap \alpha$ . A subgraph  $\vec{H}$  is *maximal* if it coincides with the initial digraph  $\vec{G}$  without one vertex  $v$  and all the edges incident to it. An *extension* of digraph  $\vec{G} = (V, \alpha)$  is the digraph  $\vec{H} = (W, \beta)$  such that, firstly,  $|W| = |V| + 1$  and, secondly, the digraph  $\vec{G}$  is includable into each maximal subgraph of the digraph  $\vec{H}$ . Consider digraphs  $\vec{G} = (V, \alpha)$  and  $\vec{H} = (W, \beta)$  such that  $V \cap W = \emptyset$ . Their *union* is the digraph  $\vec{G} \cup \vec{H} = (V \cup W, \alpha \cup \beta)$ . Now for digraphs  $\vec{G} = (V, \alpha)$  and  $\vec{H} = (W, \beta)$  such that  $V \cap W = \emptyset$  we introduce the notion of a *composition* as the digraph  $\vec{G} + \vec{H} = (V \cup W, \alpha \cup \beta \cup V \times W \cup W \times V)$ . An *isomorphism* of digraphs  $\vec{G} = (V, \alpha)$  and  $\vec{H} = (W, \beta)$  is a bijective correspondence  $\phi : V \rightarrow W$  preserving the adjacency relation, i.e.,  $((u, v) \in \alpha \leftrightarrow (\phi(u), \phi(v)) \in \beta) \quad \forall u, v \in V$ . We denote the *isomorphism relation* on digraphs  $\vec{G}$  and  $\vec{H}$  by  $\vec{G} \cong \vec{H}$ . The digraphs  $\vec{G}$  and  $\vec{H}$  are then said to be *isomorphic* [1].

A *trivial extension (TE)* of the digraph  $\vec{G} = (V, \alpha)$  is a composition  $\vec{G} + w$  of the initial digraph  $\vec{G}$  with some vertex  $w \notin V$ . Since the digraph  $\vec{G}$  trivial extension is unique up to some isomorphism, it is possible to introduce the function  $\text{TE}(\vec{G})$ . A *T-irreducible extension (TIE)* of the digraph  $\vec{G}$  is the extension of the initial graph  $\vec{G}$  being the result of the maximal possible number of edges elimination from  $\text{TE}(\vec{G})$  [2]. S. G. Kurnosova considered TIE for some cases of undirected (with symmetrical or anti-reflexive equivalence relation) graphs in [2, 3]. The authors of [3] endowed complete binary trees with TIE.

\*E-mail: alexandergavrikov1989@gmail.com.